

Presupposition Accommodation as Exception Handling

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Abstract

Van der Sandt’s algorithm for handling presupposition is based on a “presupposition as anaphora” paradigm and is expressed in the realm of Kamp’s DRT. In recent years, we have proposed a type-theoretic rebuilding of DRT that allows Montague’s semantics to be combined with discourse dynamics. Here we explore van der Sandt’s theory along the line of this formal framework. It then results that presupposition handling may be expressed in a purely Montagovian setting, and that presupposition accommodation amounts to exception handling.

1 Introduction

Montague (1970) argued that there is no essential difference between natural and mathematical languages. He developed a theory that assigns a lambda-term for each lexical item, and the meaning of a whole sentence could be obtained by composing the lambda-terms via functional application. However, his theory was limited to single sentences. De Groote (2006) extends Montague’s framework with a continuation-passing-style technique, developing a framework that is dynamic in a sense reminiscent of Dynamic Predicate Logic (Groenendijk and Stokhof, 1991).

While Montague’s semantics is based on Church’s (1940) simple type theory and has only two atomic types (ι , the type of individuals; and o , the type of propositions), de Groote (2006) adds an atomic type γ representing the type of the *environment*. For each lambda-term the continuation is what is still to be processed, and its type is $\gamma \rightarrow o$.

Since anaphoric expressions are known to be similar to presuppositional expressions (van der Sandt, 1992), it is natural to ask whether our type-theoretic framework can be extended to handle

presuppositions. The goal of this paper is to answer this question positively, at least in the case of presuppositions triggered by definite descriptions. To achieve this goal γ will not be defined simply as a list of individuals, but as a list of individuals together with their properties.

2 Background

Van der Sandt (1992) argues that presuppositions and anaphors display similar behavior: they primarily have to be bound to some antecedent previously introduced in the discourse. Therefore, they can be treated by similar mechanisms. He implements his ideas in DRT (Kamp and Reyle, 1993) in such a way that for each new sentence a provisional DRS encoding possible anaphoric elements is constructed. This provisional DRS is then merged with the main DRS, and the presuppositional anaphors are resolved in accordance with certain pragmatic constraints, so that presuppositions can be accommodated when lacking a suitable antecedent.

Geurts (1999) proposes an improvement of van der Sandt’s theory, called *the binding theory*, according to which anaphora is a kind of presupposition. Therefore, presuppositions triggered by pronouns and definite descriptions can also be accommodated: a referent is introduced with a poor descriptive content and the descriptive content can be enhanced as the discourse unfolds. Moreover, according to the presuppositional version of the *quotation theory of names* (Kneale, 1962), names (e.g. *John*) are synonymous with definite noun phrases of the form “the individual named John”. Hence, presuppositions triggered by names and by definite descriptions can be handled similarly.

De Groote’s (2006) dynamic theory provides some improvement over classical DRT. It allows the representations of sentence and discourse to be built from the lexical items in the spirit of Montague. It provides reference marker renaming for

free and may be implemented using well established techniques. We claim that Geurts’ binding theory can be incorporated into this framework, providing a fully compositional treatment of definite descriptions.

3 Presupposition in Dynamic Theory

We focus here on presuppositions triggered by definite descriptions, particularly by proper names, pronouns and possessive noun phrases.

3.1 Basic Principles

Imagine that somebody is about to tell a new story and the first sentence of this story is (1).

This story is about John. (1)

If the listener does not know John, he or she will immediately imagine a person named “John” and memorize it. In other words, the listener will *accommodate* the presuppositional content triggered by the proper name *John* in the following way: he or she will *create* a slot in the *environment*, which is some unit representing the knowledge about *John*, and put there what was just learned about *John*. Therefore, the listener will be able to refer to the created slot representing *John* as the discourse evolves. Moreover, the slot for *John* will be different from other slots, i.e. it will have some identity marker, which we call, following Karttunen (1976), *reference marker* or simply *referent*. There is a direct analogy between memory slots introduced above and Heim’s (1982; 1983) file cards: they are both aimed to store what has been learned about some individual.

Let j be the referent for *John* and assume that sentence (1) is followed by sentence (2).

John loves Mary. (2)

Mary is a new individual in the discourse and therefore *Mary* will be accommodated introducing a reference marker m exactly as it happened for *John* after the utterance of (1). The story is different for *John* now. The listener already has a representation standing for *John* in the environment, and he or she just has to turn to the corresponding slot (*select* the marker in the environment) and update the slot with the new information that John loves Mary (*bind John* from (2) to the referent j).

3.2 Proper Names

To encode, following Montague’s legacy, the observations discussed above as lambda-terms, we

first define a selection function *sel* as a function taking two arguments: a property and an environment; and returning a reference marker:

$$sel : (\iota \rightarrow o) \rightarrow \gamma \rightarrow \iota \quad (3)$$

According to Montague, proper names can be interpreted as type-raised individuals, thus the lambda-term standing for *John* in Montague’s semantics is (4), where \mathbf{j} is a constant.

$$\llbracket John \rrbracket = \lambda P.P\mathbf{j} \quad (4)$$

In the dynamic interpretation, instead of the constant \mathbf{j} we would like to have a referent corresponding to *John*. For this, we attempt to select such a referent given a property of being named John, as shown in (5).

$$\llbracket John \rrbracket = \lambda P.P(sel(\text{named “John”})) \quad (5)$$

Whether the selection of the marker for *John* succeeds depends on the current environment. Hence, instead of using Montague’s individuals (i.e. of type ι) directly, we use individuals parameterized by the environment (i.e. having type $(\gamma \rightarrow \iota)$).

Noun phrases are regarded as having type (6), which is analogous to the type for noun phrases (7) given by Montague, i.e. a noun phrase is interpreted by a lambda-term that accepts a property and returns a proposition. The only difference is that now individuals are always parameterized by an environment, and propositions are dynamic¹, i.e. they have type Ω that is defined as $\gamma \rightarrow (\gamma \rightarrow o) \rightarrow o$.

$$\llbracket NP \rrbracket = ((\gamma \rightarrow \iota) \rightarrow \Omega) \rightarrow \Omega \quad (6)$$

$$\llbracket NP \rrbracket = (\iota \rightarrow o) \rightarrow o \quad (7)$$

3.3 Pronouns

Pronouns are also presupposition triggers. It can be seen in the case of cataphora, such as, for example, in sentence (8), where in the first part of the sentence the pronoun *he* introduces an individual. Since pronouns have poorer descriptive content than proper names and they have the type of noun phrases (6), they are represented by lambda-terms that are at most as complex as the terms for proper names. The term for the pronoun *he* is shown in (9), which expresses an attempt to select a human individual having masculine gender.

When he woke up, Tom felt better. (8)

¹Analogously, dynamic predicates take two additional arguments (environment, of type γ , and continuation, of type $(\gamma \rightarrow o)$) compared to Montague’s interpretation.

$$\llbracket he \rrbracket = \lambda P.P(sel(\lambda x.human(x) \wedge masculine(x))) \quad (9)$$

If the sentence (8) is uttered in a discourse that does not provide a suitable referent, the presupposition triggered by *he* will be accommodated (as it happened for *John* in (1) and for *Mary* in (2)). The presuppositional anaphora triggered by *Tom* in the second part of the sentence could be successfully bound to the introduced referent.

3.4 Possessives

Consider the sentence (10), where we have a possessive noun phrase *John's car* triggering a presupposition that there is a car owned by John.

$$John's\ car\ is\ red. \quad (10)$$

The desired interpretation of *John's car* is shown in (11), which requires a search in the environment for a referent having the property of being a car possessed by John. The embedded presupposition is encoded via a selection function (for the inner presupposition triggered by *John*) embedded into another selection function (for the outer presupposition related to *car*).

$$\llbracket John's\ car \rrbracket = \lambda P.P(\lambda e.sel(\lambda x.car.x \wedge \mathbf{poss}\ x\ sel(\text{named "John"}))e) \quad (11)$$

However, we would like to express *John's car* compositionally in terms of its constituents. To do so, we define a term (12) taking two arguments - a noun phrase standing for a possessor and a noun standing for an object being possessed, and returning a noun phrase in form of (11). \wedge is a dynamic conjunction having type (13) and defined in (14).

$$\llbracket 's \rrbracket = \lambda YX.\lambda P.P(SEL(\lambda x.((Xx) \wedge Y(\llbracket poss \rrbracket x)))) \quad (12)$$

$$\wedge : \Omega \rightarrow (\Omega \rightarrow \Omega) \quad (13)$$

$$A \wedge B = \lambda e\phi.Ae(\lambda e.Be\phi) \quad (14)$$

The term $\llbracket poss \rrbracket$ in (12) is a usual dynamic two-arguments predicate, its lambda-term is shown in (15). *SEL* is a higher-order selection function. It has the same designation as (3), with the only difference that it functions on the level of dynamic propositions. Thus, the type of *SEL* is (16) and it is analogous to the type of *sel* spelled in (3). Moreover, *SEL* is defined via *sel*, and the corresponding lambda-term is presented in (17).

$$\llbracket poss \rrbracket = \lambda xy.\lambda e\phi.\mathbf{poss}(xe)(ye) \wedge \phi e \quad (15)$$

$$SEL : ((\gamma \rightarrow \iota) \rightarrow \Omega) \rightarrow \gamma \rightarrow (\gamma \rightarrow \iota) \quad (16)$$

$$SEL = \lambda Pe.sel(\lambda x.P(\lambda e.x)e(\lambda e.\top))e \quad (17)$$

$$\llbracket car \rrbracket = \lambda x.\lambda e\phi.\mathbf{car}(xe) \wedge \phi e \quad (18)$$

If we apply the term $\llbracket 's \rrbracket$ to the term (5) for *John* and the term (18) for *car*, which is just a dynamic unary predicate, we will get the desired result (11).

3.5 Implicit Referents

Sometimes an anaphora wants to be bound, even though no referent was introduced explicitly, as in (19). Already after the first sentence, a listener will learn that John has a wife, i.e. introduce a new referent. The presuppositional anaphora triggered by the possessive noun phrase *his wife* in the second sentence will be bound to this referent.

$$John\ is\ married.\ His\ wife\ is\ beautiful. \quad (19)$$

This case can be accounted with the lexical interpretation in (20) for *being married*, which is defined by a two-arguments relation **is_married**. The first argument of the relation is the argument x being passed to the lexical interpretation. The second argument is an individual selected from the environment given the property of being either the wife or the husband of x .

$$\llbracket is_married \rrbracket = \lambda x.\lambda e\phi.\mathbf{is_married}(xe)(sel(\lambda y.(\mathbf{wife}(y,x) \vee \mathbf{husband}(y,x)))e) \wedge \phi e \quad (20)$$

3.6 Discourse Update

A discourse is updated by appending the next sentence, as shown in equation (21). A sentence is defined as a term having the type of a dynamic proposition, i.e. its type is (22), while a discourse is defined as a term having the type of a dynamic proposition evaluated over the environment, i.e its type is (23). A discourse \mathbb{D} updated with a sentence \mathbb{S} results in a term having type (23), thus it has one parameter ϕ of type $(\gamma \rightarrow o)$. The body must be a term, of type o , contributed by \mathbb{D} . \mathbb{D} itself is a term of type (23). Therefore, it must be given a continuation as an argument constructed with \mathbb{S} and its continuation.

$$\mathbb{D} \cup \mathbb{S} = \lambda \phi.\mathbb{D}(\lambda e.\mathbb{S}e\phi) \quad (21)$$

$$\llbracket S \rrbracket = \Omega = \gamma \rightarrow (\gamma \rightarrow o) \rightarrow o \quad (22)$$

$$\llbracket D \rrbracket = (\gamma \rightarrow o) \rightarrow o \quad (23)$$

However, during the computation of $\lambda \phi.\mathbb{D}(\lambda e.\mathbb{S}e\phi)$ one of the selection functions can raise an exception containing a message that a referent having some property Q was not found in the environment. The exception will be caught and the property will be returned to the exception

handler. The handler will have to introduce a referent having the property Q into the representation of the discourse, add this referent to the environment, and call the update function passing to it the amended interpretation of the discourse and the sentence \mathbb{S} as parameters. This can be encoded using an exception handling mechanism as shown in (24) for global accommodation. Note that the definition of discourse update is recursive.

$$\begin{aligned} \mathbb{D} \cup \mathbb{S} &= \lambda\phi. \mathbb{D}(\lambda e. \mathbb{S}e\phi) \\ &\text{handle } (\text{fail } Q) \text{ with} \\ &\lambda\phi. \mathbb{D}(\lambda e. \exists x. (Qx) \wedge \phi((x, Qx) :: e)) \cup \mathbb{S} \end{aligned} \quad (24)$$

The environment is defined as a list of pairs “referent \times proposition” (25). The two-place list constructor $::$ appends a referent together with the corresponding propositions into the environment, therefore it has the type shown in (26).

$$\gamma = \text{list of } (\iota \times o) \quad (25)$$

$$:: : (\iota \times o) \rightarrow \gamma \rightarrow \gamma \quad (26)$$

The selection function sel can implement any anaphora resolution algorithm, and hence our framework is not confined to any of them.

Considering that the lambda-term for *Mary* is similar to (5) and the lambda-term for the transitive verb *love* is (27), the interpretation for the sentence (2) after beta-reductions will be (28).

$$\llbracket \text{love} \rrbracket = \lambda YX.X(\lambda x.Y(\lambda y.(\lambda e\phi.\text{love}(xe)(ye) \wedge \phi e))) \quad (27)$$

$$\begin{aligned} \mathbb{S}_2 &= \llbracket \text{love} \rrbracket \llbracket \text{John} \rrbracket \llbracket \text{Mary} \rrbracket \rightarrow_{\beta}^* \\ &\lambda e\phi. (\text{love}(sel(\text{named } \text{“John”})e) \\ &\quad (sel(\text{named } \text{“Mary”})e)) \wedge \phi e \end{aligned} \quad (28)$$

After the sentence (1), the lambda-term representing discourse will be (29).

$$\begin{aligned} \mathbb{D}_1 &= \lambda\phi. \exists y. (\text{story } y) \wedge \\ &\quad \exists j. (\text{named } \text{“John” } j) \wedge \\ &\quad \text{about } (y, j) \wedge \\ &\quad \phi((y, \text{story } y) :: (j, \text{named } \text{“John” } j)) \end{aligned} \quad (29)$$

After the sentence (2), the lambda-term \mathbb{D}_1 in (29) will have to be updated with the term \mathbb{S}_2 in (28) as it is defined by the function (24). Since we have a referent for *John* in the environment of \mathbb{D}_1 , it will be successfully selected and *John* from \mathbb{S}_2 will get bound to it. However, there will be a failure for *Mary*, particularly on the property (named “Mary”) since there is no corresponding referent in \mathbb{D}_1 yet. The failure will be handled by accommodating *Mary* and introducing the sentence \mathbb{S}_2 into the amended interpretation of the discourse, which results in the term shown in (30).

$$\begin{aligned} \mathbb{D}_2 &= \mathbb{D}_1 \cup \mathbb{S}_2 = \lambda\phi. \exists y. (\text{story } y) \wedge \\ &\quad \exists j. (\text{named } \text{“John” } j) \wedge \\ &\quad \text{about } (y, j) \wedge \\ &\quad \exists m. (\text{named } \text{“Mary” } j) \wedge \\ &\quad \text{love } (j, m) \wedge \\ &\quad \phi((m, \text{named } \text{“Mary” } m) :: \\ &\quad (y, \text{story } y) :: \\ &\quad (j, \text{named } \text{“John” } j)) \end{aligned} \quad (30)$$

4 Conclusions

We showed that de Groote’s (2006) dynamic framework can be applied to presuppositions triggered by definite descriptions, such as proper names, possessive noun phrases and pronouns; and that the exception handling mechanisms offer a proper way of modeling the dynamics of presupposition. Other presuppositional expressions, such as, for example, factives and aspectual verbs, will require more technicalities. Nevertheless, we believe that the approach can be extended to encompass a general theory of presupposition and we intend to address this in future work.

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